

Monoids revisited

Comm. monoids: M with
in a cat. C

$$M \times M \xrightarrow{\mu} M \quad \text{binary operation}$$

$$\ast \xrightarrow{e} M \quad \text{nullary operation}$$

+ axioms

The axioms ensure that for any finite set I , there is a well-defined multiplication map $\mu_I : M^I \rightarrow M$.

$$Y = \coprod_{i \in I} Y_i \rightsquigarrow M^Y \cong \prod_{i \in I} M^{Y_i} \xrightarrow{\prod \mu_{Y_i}} \prod_{i \in I} M.$$

$\downarrow \mu_Y$ ↗ $\prod_{i \in I} M$

$$Y \xrightarrow{f} I \quad Y_i = f^{-1}(i)$$

for any $f: Y \rightarrow I$, we have $\mu_f: M^Y \rightarrow M^I$

$$K \xrightarrow{g} Y \xrightarrow{f} I \rightsquigarrow \mu_f \circ \mu_g = \mu_{f \circ g}$$

The comm. monoid structure on M defines a functor $\mu: \text{Fin} \rightarrow C$

$$I \longmapsto M^I$$

$$(f: Y \rightarrow I) \mapsto \mu_f.$$

Q: Can we recover M from the functor μ ?

Not quite...

We also need isomorphisms $\mu(I) \cong \mu(\ast)^I$

$$\downarrow \pi_i \quad i \in I$$

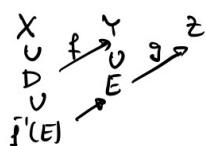
$$\mu(\ast)$$

We need to add morphisms to Fin :

$$\begin{array}{c} \text{Fin} \subset \text{Fin}' \\ \text{(Segal)} \\ \Downarrow \\ \text{Fin}'^\text{op} \end{array}$$

Fin' : objects are finite sets
morphisms are partially defined maps
• A part. defined map from X to Y
is a pair (D, f) where $D \subset X$
and $f: D \rightarrow Y$

Composition:



Remark: $\text{Fin}' \cong \text{Fin}_*$

$$\begin{array}{ccc} A\circ & & A_+ = A \cup \{\star\} \\ \downarrow f^{-1}(B) & \leftarrow & \downarrow f \\ B & \xrightarrow{f} & B_+ \end{array}$$

Given $\mathbb{J} \supseteq D \xrightarrow{f} I$, we get $M^D \xrightarrow{\text{res}} M^I \xrightarrow{\mu_I} M^I$

So a commutative monoid M in C defines a functor

$$\text{Fin}' \rightarrow C, \quad I \mapsto M^I.$$

$$i \in I \in \text{Fin}, \quad \rho_i : I \supseteq \{i\} \rightarrow * \quad \text{"Segal map"}$$

$$\rightsquigarrow M^I \xrightarrow{\pi_i} M$$

Prop Let C be a category with finite products. There is a fully faithful functor

$$\begin{aligned} \text{CMon}(C) &\hookrightarrow \text{Fun}(\text{Fin}', C) \\ M &\longmapsto (I \mapsto M^I) \end{aligned}$$

A functor $X : \text{Fin}' \rightarrow C$ is in the essential image iff for every $I \in \text{Fin}$,

$$X(I) \xrightarrow{(f_i)_{i \in I}} X(*)^I \quad \text{"Segal condition"}$$

is an isomorphism.

Pf. Exercise.

Remark One can show that there is an equivalence of 2-categories

$$\begin{aligned} \text{SymMonCat} &\stackrel{\text{Sym}}{\simeq} \text{Fun}(\text{Fin}', \text{Cat}) \subset \text{Fun}(\text{Fin}', \text{Cat}) \\ &\text{full subcategory on } X : \text{Fin}' \rightarrow \text{Cat} \\ &\text{st: } X(I) \xrightarrow{(f_i)_{i \in I}} X(*)^I \text{ is an equivalence.} \end{aligned}$$

General monoids $M \in \text{Mon}(C)$

To define $\mu_I : M^I \rightarrow M$, we need an ordering on I .

Let Ord be the category of finite ordered sets,

As before, M defines $\mu : \text{Ord} \rightarrow C$

$$\begin{array}{c} I \mapsto M^I \\ (\mathbb{J} \xrightarrow{f} I) \mapsto \mu_f : M^{\mathbb{J}} \xrightarrow{\downarrow f^{-1}(i)} M^I \xrightarrow{\mu} M \end{array}$$

Let $\text{Ord}_{\pm\infty}$ be the subcategory of Ord with objects $I_{\pm\infty} = I \cup \{\pm\infty\}$
and morphisms satisfy $f(-\infty) = -\infty$
 $f(+\infty) = +\infty$.

$$\text{Ord} \hookrightarrow \text{Ord}_{\pm\infty} \quad \left(\begin{array}{l} \text{analogous to: } \text{Fin} \hookrightarrow \text{Fin}_* \\ I \hookrightarrow I_{\pm\infty} \end{array} \right)$$

$$\begin{array}{ccc} -\infty & \cdot & [\cdot & \cdot & \cdot] & +\infty \\ \searrow & & \downarrow & & \downarrow & \\ -\infty & \cdot & [\cdot & \cdot & \cdot] & +\infty \end{array}$$

Segal maps: $i \in I \in \text{Ord}$, $f_i: I_{\pm\infty} \rightarrow *_{\pm\infty}$

$$j \mapsto \begin{cases} * & \text{if } j=i \\ -\infty & \text{if } j < i \\ +\infty & \text{if } j > i \end{cases}$$

Prop Let C be a category with finite products. There is a fully faithful functor $\text{Mon}(C) \hookrightarrow \text{Fun}(\text{Ord}_{\pm\infty}, C)$

A functor $X: \text{Ord}_{\pm\infty} \rightarrow C$ is in the resulting image iff for every $I \in \text{Ord}$, $X(I_{\pm\infty}) \xrightarrow{(p_i)_{i \in I}} X(*_{\pm\infty})^I$ is an isomorphism.

Remark One can show that there is an equivalence of 2-categories

$$\text{MonCat} \simeq \text{Fun}_{\text{Segd}}(\text{Ord}_{\pm\infty}, \text{Cat})$$

Exercise there is an equivalence $\text{Ord}_{\pm\infty} \simeq \Delta^{\text{op}}$

$$\{1, \dots, n\}_{\pm\infty} \mapsto [n]$$

Under this equivalence, the Segal maps are $p_i: [1] \rightarrow [n] \quad (1 \leq i \leq n)$

$$\begin{array}{c} 0 \mapsto i-1 \\ 1 \mapsto i \end{array}$$

Remark 1) The inclusion $C\text{Mon}(C) \hookrightarrow \text{Mon}(C)$ is induced by the functor

$$\text{Ord}_{\pm\infty} \rightarrow \text{Fin}_* \text{ sending } I_{\pm\infty} \text{ to } I_+.$$

2) For $C = \text{Set}$, this is the nerve construction:

$$\text{Mon} \hookrightarrow \text{Cat} \xrightarrow{N} \text{Set}$$

$$M \longmapsto (* \subseteq M \subseteq M \times M \dots)$$